

# COMP90042 Web Search & Text Analysis

Workshop Week 3

---

Zenan Zhai

March 19, 2019

University of Melbourne

## Indexing

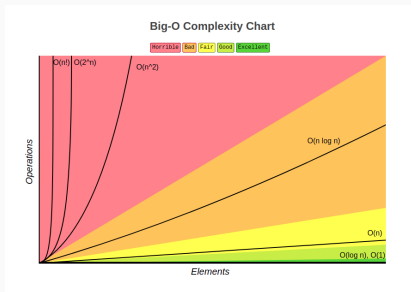
- Data Structure
  - Document-Term Matrix
  - Inverted Index
- Compression
  - Variable Byte Compression
  - OptPFor Delta Compression
- Index Construction
  - Invert Batch Indexing
  - Auxiliary Indexing
  - Logarithmic Indexing

## Search

- Vector Space Models
  - TF-IDF
  - BM25
- Efficient Query Processing
  - Operation GEQ
  - WAND
- Query Completion
  - Prefix Trie
  - Range Maximum Query
- Query Expansion
  - Relevance Feedback
  - Semantic-Based Methods
- Phrase Search
  - Inverted Index + Positional Information
  - Suffix Array
- Evaluation and Re-rank

# Warm Up - Complexity

## Time Complexity



Big O cheat sheet

## Notation

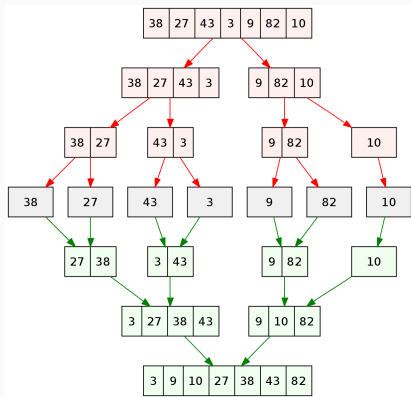
- $T(n) = O(f(n)) \Leftrightarrow \exists c, n_0, \forall n > n_0, T(n) \leq c \cdot O(f(n))$
- $T(n) = \Omega(f(n)) \Leftrightarrow \exists c, n_0, \forall n > n_0, T(n) \geq c \cdot \Omega(f(n))$
- $T(n) = \Theta(f(n)) \Leftrightarrow \exists c_1, c_2, n_0, \forall n > n_0,$   
 $c_1 \cdot \Theta(f(n)) \leq T(n) \leq c_2 \cdot \Theta(f(n))$

# Warm Up - Complexity

## Space Complexity

- Amount of auxiliary space the algorithm need in the function of input size.

## Example - Merge Sort



*Fig from Wikipedia.*

- Index Compression
  - Variable Byte Compression
  - OptPFor Delta Compression
  
- Index Construction
  - Invert Batch Indexing
  - Auxiliary Indexing
  - Logarithmic Indexing
  
- Efficient Query Processing
  - Operation GEQ
  - WAND

# Variable Byte Compression

*1Byte = 8bits*

Biggest integer can be stored in 64/32 bits.

- Max. 64-bits unsigned integer  
 $2^{64} - 1 = 18,446,744,073,709,551,615$
- Max. 32-bits unsigned integer  
 $2^{32} - 1 = 4,294,967,295$

In practice, 95% of integers in posting lists are  $< 128$   
(Can be stored in 7 bits).

# Variable Byte Compression

## Compression of integer 315700

- Divide the integer by 128 ( $2^7$ ).
- Add "indicator bit" to the residual.  
1 indicates end of encoded number.

315700  $\Rightarrow$  0010011|0100010|0110100

Indicator	Remainder	Decimal
0	0110100	52
0	0100010	34
1	0010011	147

## Decompression of integer 315700

- Multiply the remainder by  $2^{7 \times (i-1)}$  for the  $i^{\text{th}}$  chunk.
- Take sum of all chunks.

$$52 \times 2^0 + 34 \times 2^7 + (147 - 128) \times 2^{14} = 315700$$



## More flexible chunk size $b$ than VByte compression

- Encode most numbers in a block by  $b$  bits
- Leave exception unchanged.
- Record number and positions of exceptions in header.

Encode [1 4 7 2 4 5 123 48] with  $b = 3$

Header	Content	Exceptions
$b=3, \#e=2, epos=[6, 7]$	[ 1, 4, 7, 2, 4, 5 ]	[123, 48]

Parameter  $b$  need to be tuned to get optimal performance.

- Index Compression
  - Variable Byte Compression
  - OptPFor Delta Compression
- Index Construction
  - Static Construction
  - Auxiliary Indexing
  - Logarithmic Indexing
- Efficient Query Processing
  - Operation GEQ
  - WAND

# Static Construction

Traverse once to get global vocabulary.

{1:Apple, 2:Pear, 3:Banana, 4:Grape}

Split document collection into batches and compute inverted indexes in parallel, then merge terms with same ID.

Apple  $\rightarrow d_1, d_2$

Pear  $\rightarrow d_2$

Banana  $\rightarrow d_1$

+

Apple  $\rightarrow d_3$

Pear  $\rightarrow d_3, d_4$

Grape  $\rightarrow d_4$

$\Downarrow$

Apple  $\rightarrow d_1, d_2, d_3$

Pear  $\rightarrow d_2, d_3, d_4$

Banana  $\rightarrow d_1$

Grape  $\rightarrow d_4$

# Auxiliary Indexing

## Combination of static and incremental indexing

- 1 static index on disk, 1 incremental index on memory.
- Merge when incremental index is too big.
- Query both indexes and merge results

Note that merge requires number of IOs equals to the size of the posting lists

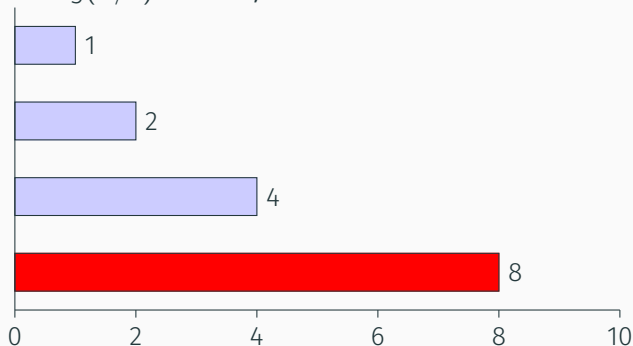
Suppose now we have  $N$  postings on disk and we merged the postings when it reaches size  $n$ .

$$T(N, n) = \sum_{i=1}^{\frac{N}{n}} i \times n = \frac{(1 + \frac{N}{n})}{2} n \times \frac{N}{n} = \frac{N}{2} + \frac{N^2}{2n} (n < N)$$

$$T(N, n) = O\left(\frac{N^2}{n}\right)$$

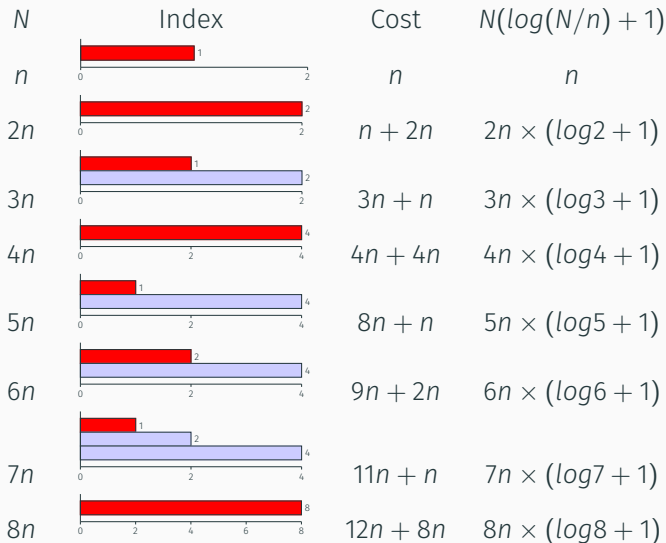
# Logarithmic Indexing

Use  $\log(N/n)$  indexes, each level  $i$  has size  $2^i \times n$ .

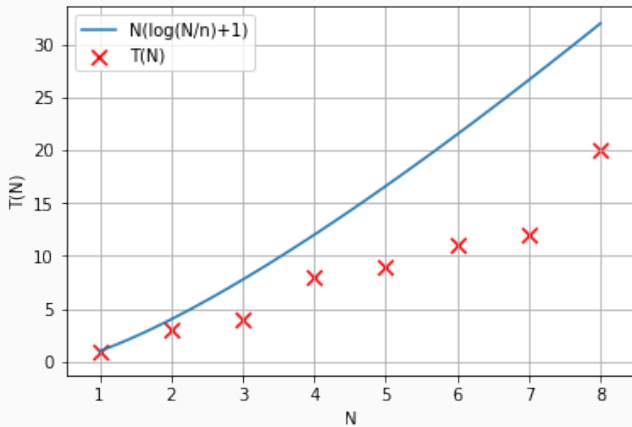


Consider total number of merges needed when  $N = 8n$ .

# Logarithmic Indexing



# Logarithmic Indexing



- Index Compression
  - Variable Byte Compression
  - OptPFor Delta Compression
  
- Index Construction
  - Static Construction
  - Auxiliary Indexing
  - Logarithmic Indexing
  
- Efficient Query Processing
  - Operation GEQ
  - WAND



GEQ - Greater or equal to x.

2, 4, 7, 9, 13, 20, 34, 54, 67, 87, 96, 103

Binary search in sorted list.

- Time complexity ?
- Space complexity ?

Compressed posting lists with sample value

[2, 4, 7, 9]    [13, 20, 34, 54]    [67, 87, 96, 103]

max=9

max=54

max=103

$$W_{d,t} = \frac{(k_1 + 1)tf_{d,t}}{k_1((1 - b) + b(\frac{L}{L_{avg}})) + tf_{d,t}} \times \log \frac{N - df_t + 0.5}{df_t + 0.5} \times \frac{(k_3 + 1)tf_{Q,t}}{k_3 + tf_{Q,t}}$$

- $T(n)$  for all words in  $Q$  with 1 document

$$T(\text{score}(Q, d)) = \sum_{q \in Q} W_{d,q} = O(|Q| + |Q|) = O(|Q|)$$

- $T(n)$  for 1 word  $Q$  with all document

$$T(\text{score}(q, D)) = \sum_{d \in D} W_{d,q} = O(|D| + 1) = O(|D|)$$

- $T(n)$  for all words in  $Q$  with all document

$$T(\text{score}(Q, D)) = \sum_{q \in Q} \sum_{d \in D} W_{d,q} = O(|Q| \times |D| + |Q|) = O(|Q| \times |D|)$$

## Concepts

- Top- $k$  retrieval
- Maximum Contribution

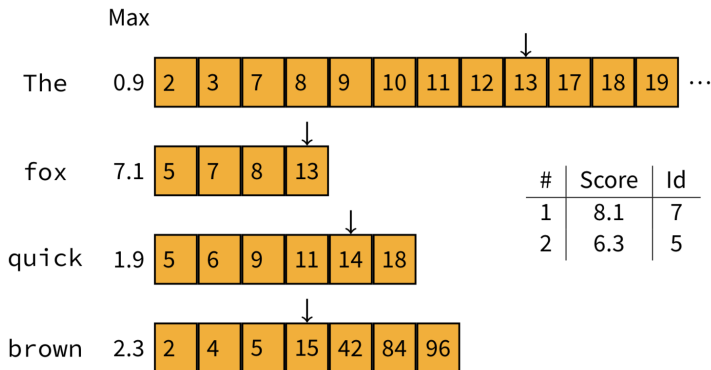
## Questions

- Why do we want to sort lists by their current pivot?
- Why do we still need to compute similarity score if Max. contribution of a doc is greater than Min. score in top- $k$  list?
- How does operation GEQ helps WAND?

## Benefits and Restrictions

# Exercise - What does the WAND do?

Query Q: The quick brown fox with  $k = 2$

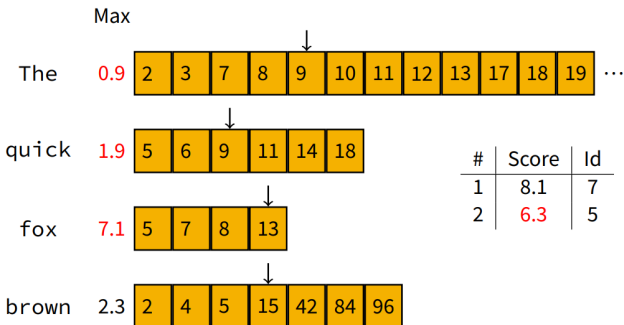


# Exercise - What does the WAND do?

## WAND - Example - Fast Forward

Query Q: The quick brown fox with  $k = 2$

What is the next document that has to be evaluated?



**13**, as  $0.9 + 1.9 + 7.1 > 6.3!$